

Equivalence Relations

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Property

1. $x \in [x]$
2. $x R y$ iff $[x] = [y]$
3. $[x] = [y]$ or $[x] \cap [y] = \emptyset$

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 $x \in [x]$.

2. " \Rightarrow " If $x R y$, let $w \in [x]$.
Then $w R x$ and hence $w R y$ (since R is transitive).
We have $w \in [y]$. Therefore, $[x] \subseteq [y]$.
If $t \in [y]$, then $t R y$.
As $y R x$ (since R is symmetric) we have $t R x$.
Hence $t \in [x]$. Therefore, $[y] \subseteq [x]$.

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$$\therefore [x] = [y].$$

" \Leftarrow " Suppose $[x] = [y]$. Since $x \in [x]$,
 $x \in [y]$ and hence $x R y$.

3. If $x R y$, from part 2, $[x] = [y]$.
If $x \not R y$, let $[x] \cap [y] \neq \emptyset$.

There exists v such that $v \in [x]$ and $v \in [y]$.

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Hence $x R v$ and $v R y$, which implies
that $x R y$, contradicting the assumption
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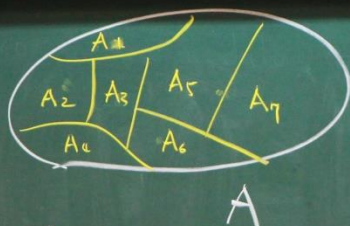
Def A partition of a set A is a family
 $\{A_i : i \in I\}$ of nonempty subsets of A
such that

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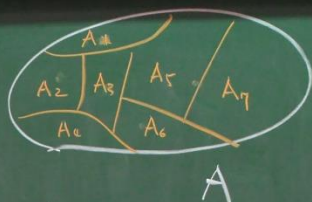
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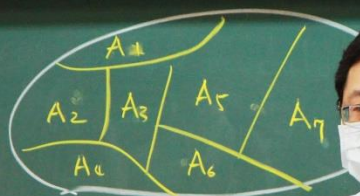
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$$2. A_i \cap A_j = \emptyset \text{ for all } i \neq j, i, j \in I$$



Property If R is an equivalence relation on A then the distinct equivalence classes (with R) form a partition of A .

- $\Rightarrow x R x$.
2. Symmetric: If x and y belong to the same subset, then y and x belong to the same subset.
3. Transitive: If x and $y \in A_i$ and y and $z \in A_j$ then $A_i = A_j$ (since $\{A_i : i \in I\}$ is a partition). Hence, x and $z \in A_i$.

Proof This follows directly from parts 1 and 3 of the previous property.

Property Let $\{A_i : i \in I\}$ form a partition of A . Define $x R y$ iff x and y belong to the same subset of partition. Then R is an equivalence relation.

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Property Let $\{A_i : i \in I\}$ form a partition of A .
Define $x R y$ iff x and y belong to the same subset
of partition. Then R is an equivalence relation.

Since $x \in U$ and $v \in K y$, which implies that $x R y$, contradicting the assumption that $x \not R y$. Therefore, we must have $[x] \cap [y] = \emptyset$.

Remark For any set A , there is a one-to-one correspondence (bijection) between the set of equivalence relations on A and the set of partitions of A .

Example Congruence modulo m on \mathbb{Z} :
Each integer is congruent modulo m to one and only one of $0, 1, 2, \dots, m-1$.

$$\begin{aligned}
 [0] &= \{ \dots, -2m, -m, 0, m, 2m, \dots \} \\
 [1] &= \{ \dots, -2m+1, -m+1, 1, m+1, 2m+1, \dots \} \\
 [2] &= \{ \dots, -2m+2, -m+2, 2, m+2, 2m+2, \dots \} \\
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Partial Orders

Def A relation R on a set is a partial order (or a partial ordering relation) if the following conditions hold:

1. $a R a$ for all $a \in A$ (reflexive)
2. $a R b$ and $b R a \Rightarrow a = b$ (antisymmetric)
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Example $A = \mathbb{N}$

$$a R b \Leftrightarrow a | b$$

Is " $|$ " a partial order?

1. Reflexive: $a | a$ for all $a \in \mathbb{N}$.
2. Antisymmetric: If $a | b$ and $b | a$, then

b/a and a/b are both positive integers
and $(b/a)^{-1} = a/b \therefore a/b = 1$
 $\Rightarrow a=b$.

3. **Transitive:** If $a|b$ and $b|c$, then
 $b = ak_1$ and $c = bk_2$ for some $k_1, k_2 \in \mathbb{N}$
Hence $c = a(k_1 k_2)$, implying that $a|c$.
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Example Consider the power set $\mathcal{P}(S)$ of a set S .

Then " \subseteq " is a partial order on $\mathcal{P}(S)$.

1. Reflexive: $A \subseteq A$ for all subset A of S .
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Suppose the partial order is " \leq ".

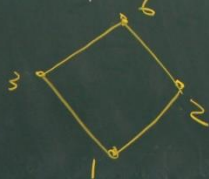
Each element of A is represented by a point,
and we connect a to b (with a below b)
with a line iff $a < b$, i.e. $a \leq b$ and $a \neq b$,
and there is no $c \in A$ such that $a < c < b$.

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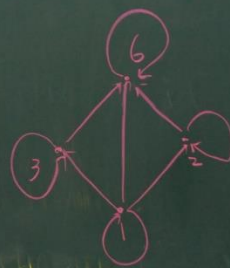
Example $A = \{a \in \mathbb{N} : a \mid 6\}$

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partial order: \mid



Hasse diagram



digraph

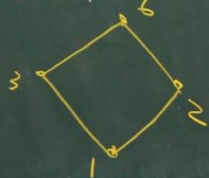
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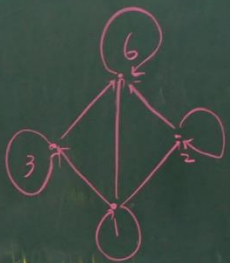
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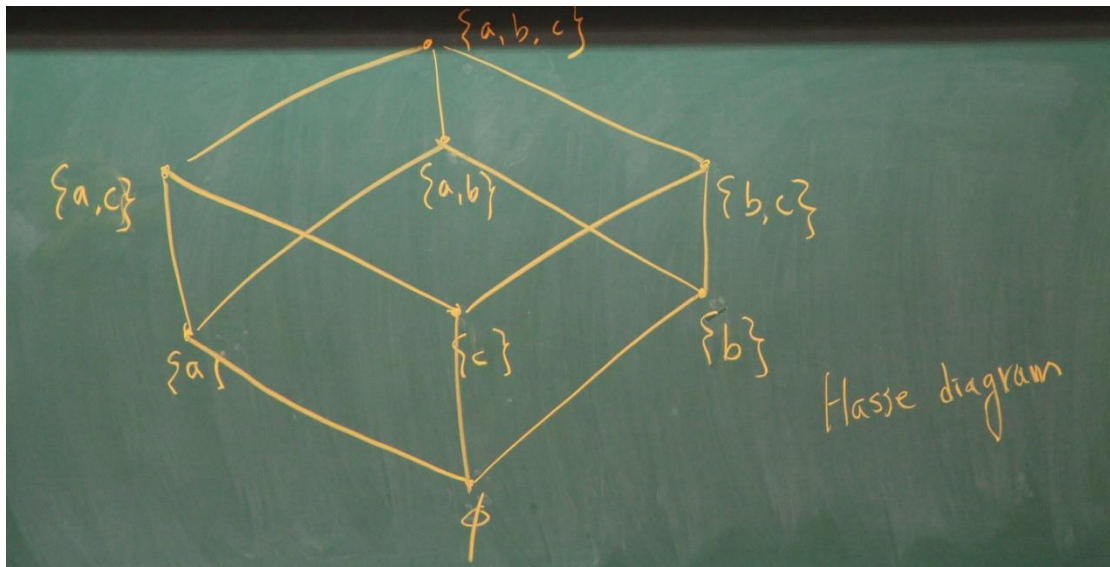


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partial order: \subseteq



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